Block cipher

Properties:
- Deterministic
- Without the key plaintext cannot be found
- Valid plaintext-ciphertext pairs do not leak the key
- Diffusion & Confusion

AES – Advanced Encryption Standard (NIST 2001)
- 16-byte (128-bit) block size
- key sizes – 128/192/256 bits
Electronic Codebook (ECB) mode

Key → Plaintext → block cipher encryption → Ciphertext

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Initialization Vector (IV)

- On encryption XOR plaintext with random IV
- IV must be sent along with the ciphertext
- IV does not have to be secret
- On decryption XOR plaintext with IV

The problem? The ciphertext is two times larger than the plaintext
Cipher Block Chaining (CBC) mode

- Serial encryption
- Parallel decryption
- What about integrity (malleability)?
Plaintext padding

- Random padding:
  1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a XX YY ZZ

- Zero padding:
  1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 00 00 00

- ANSI X.923:
  1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 00 00 03

- ISO/IEC 10126:
  1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a F1 A6 03

- ISO/IEC 7816-4:
  1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 80 00 00

- PKCS#5/PKCS#7:
  1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 03 03 03
  1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 1a 05 05 05 05 05

Padding is added even if the plaintext occupies a full block
Counter (CTR) mode

- Block cipher-based PRNG
- Turns block cipher into a stream cipher
- Nonce must never be reused!
- Block ciphers vs stream ciphers
Galois/Counter Mode (GCM)

- Authentication tag is 16 bytes (128 bits) long
- Standardized for 16-byte block size
- Nonce 12 bytes, counter 4 bytes (max 64GB of data)
- Authenticated Encryption with Associated Data (AEAD)
Disk encryption

• Encrypt the whole disk using CBC?
  • Each sector is encrypted separately
  • Use sector number as IV

• How to provide integrity?
  • MAC is not an option
  • XTS mode

• Password change without disk reencryption
  • Random master key is used to encrypt data
  • Master key is stored encrypted in disk header
  • User’s password decrypts master key
  • Enterprise solutions have several encryptions of master key
Password-based encryption

Deriving strong (128-bit) keys from short, low entropy passwords?

- Use hash of the password
  - Distributes entropy over all 128 bits
- Use salt as an addition to password
  - Prevents multi-target attacks
- Use iterated hash to slow down brute-force
  - Adds arbitrary number of operations to brute-force
Password-Based Key Derivation Function 2 (PBKDF2)

key = PBKDF2(PRF, Password, Salt, iter, kLen)

- PRF – pseudorandom function (e.g., HMAC-MD5, HMAC-SHA1)
- Password – password entered by the user
- Salt – random cryptographic salt
  - Recommended at least 64 bits
- iter – number of iterations desired
  - Recommended at least 1’000 iterations (increases the security level by 10 bits)
  - NIST recommends 10’000’000 for critical keys (increases the security level by 23 bits)
  - Ubuntu disk encryption selects iteration count that corresponds to 1 second
- kLen – desired length of the derived key
  - WPA2 uses key = PBKDF2(HMAC-SHA1, passphrase, ssid, 4096, 256)
  - Truecrypt uses PBKDF2 with 2000 iterations

New schemes should use scrypt / Argon2id
Task: Password-based file encryption

Implement a utility that encrypts and decrypts files using a password:

```
$ ./aes.py
Usage:
-encrypt <plaintextfile> <ciphertextfile>
-decrypt <ciphertextfile> <plaintextfile>
```

```
$ ./aes.py -encrypt plain.txt plain.txt.enc
[+] Benchmark: 806972 PBKDF2 iterations in 1 second
[?] Enter password: asd
```

```
$ ./aes.py -decrypt plain.txt.enc plain.txt
[?] Enter password: asd
```

- The integrity of the ciphertext is provided using HMAC
- Parameters are ASN.1 DER-encoded and stored as a header of the ciphertext file
Task: Password-based file encryption

```plaintext
$ dumpasn1 plain.txt.enc
0 102: SEQUENCE {
  2 18:   SEQUENCE {
    4 8:     OCTET STRING 59 9C B8 63 65 82 FA 39
    14 3:    INTEGER 806972
    19 1:    INTEGER 48
      :   }
  22 29:   SEQUENCE {
    24 9:     OBJECT IDENTIFIER aes128-CBC (2 16 840 1 101 3 4 1 2)
    35 16:    OCTET STRING 7A AD 60 A6 50 F2 42 49 9E 7A 34 70 39 37 CF C6
      :   }
  53 49:   SEQUENCE {
    55 13:    SEQUENCE {
      57 9:      OBJECT IDENTIFIER sha-256 (2 16 840 1 101 3 4 2 1)
      68 0:      NULL
      :   }
    70 32:    OCTET STRING
      :     FC 71 A6 AC D5 35 52 1C C7 BF 13 15 AB B8 FA 85
      :     87 81 1E C9 8F 3F 83 91 OB D7 B3 86 A2 ED 8B C3
      :   }
      :   }
  Warning: Further data follows ASN.1 data at position 104.

EncInfo ::= SEQUENCE {
  kdfInfo pbkdf2params,
  cipherInfo aesInfo,
  hmacInfo DigestInfo
}

pbkdf2params ::= SEQUENCE {
  salt OCTET STRING,
  iterationCount INTEGER (1..MAX),
  keyLength INTEGER (1..MAX)
}
aesInfo ::= SEQUENCE {
  algorithm OBJECT IDENTIFIER,
  iv OCTET STRING OPTIONAL,
}
```

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Task: Test cases

$ echo -n "hello world" > plain
$ ./aes.py -encrypt plain plain.enc
[+] Benchmark: 818397 PBKDF2 iterations in 1 second
[?] Enter password: asd
$ ./aes.py -decrypt plain.enc plain.new
[?] Enter password: asd
$ hexdump -C plain.new
00000000  68 65 6c 6c 6f 20 77 6f 72 6c 64          |hello world|
0000000b

$ echo -e -n "hello world \x01\x01\x02\x02" > plain
$ ./aes.py -encrypt plain plain.enc
[+] Benchmark: 816193 PBKDF2 iterations in 1 second
[?] Enter password: asd
$ ./aes.py -decrypt plain.enc plain.new
[?] Enter password: asd
$ hexdump -C plain.new
00000000  68 65 6c 6c 6f 20 77 6f 72 6c 64 20 01 01 02 02 |hello world ....|
00000010

$ ./aes.py -decrypt plain.enc plain.new
[?] Enter password: asdd
[-] HMAC verification failure: wrong password or modified ciphertext!

$ openssl dgst -sha256 big
SHA256(big)= 066090dceece702a28c9ab08677bd91f7e53fa6b5a69d1d4c3e9a2d556e4cee
Task: Password-based file encryption

• Slow down password brute-force attacks to 1 try/second
  • Benchmark the time required for 10 000 iterations
  • Extrapolate the iteration count to 1 second
    
    ```python
    start = time.time()
    ...
    stop = time.time()
    took = stop-start
    ```

• Use PBKDF2 with HMAC-SHA1 to derive 48 bytes:

    ```python
    pbkdf2_hmac('sha1', password, salt, iter, 48)
    ```

  • Use the first 16 bytes as AES-128 key
  • The next 32 bytes as HMAC-SHA256 key

• Generate IV (16 bytes) and salt (8 bytes) randomly

• Implement CBC mode using pure AES-128 (ECB mode)

    ```python
    cipher = AES.new(key_aes, AES.MODE_ECB)
    cipher.encrypt(strxor(plaintext_block, iv_current))
    strxor(cipher.decrypt(ciphertext_block), iv_current)
    ```

• Use PKCS#5 padding
Task: Password-based file encryption

- The entire plaintext/ciphertext can be stored and processed in memory
- Calculate the MAC on the ciphertext
  - Prepend IV to ciphertext when calculating MAC
  - Use `hmac.compare_digest(mac_calculated, mac_from_der)`
Side channel attacks

```python
def authenticate_admin(submitted_password):
    hardcoded_password = 'qwerty'
    if submitted_password == hardcoded_password:
        return 1 # access granted
    return 0 # access denied
```

- Function is vulnerable to a timing attack (comparison stops after the first incorrect byte):
  - password 'aaaaaa' – 1ms
  - password 'baaaaa' – 1ms
  - password 'qaaaaa' – 2ms
  - password 'qwaaaa' – 3ms
  - password 'qweaaa' – 4ms
  - password 'qweara' – 5ms

- Exploitable over low-latency networks
- Adding `sleep(random())` before `return` will not help
- Constant-time string comparison function needed

Your `hmac.py` solution (homework 03) is vulnerable to a timing attack

```python
hmac.compare_digest(mac_calculated, mac_from_der)
```
Questions

• How does a block cipher work (e.g., takes as input, returns)?

• What happens to the ciphertext if a single plaintext or key bit is changed?

• Why is encrypting every block of a file independently not secure?

• Why do we apply an initialization vector (IV) to plaintext?

• How can integrity be provided for ciphertext?

• When should a stream cipher be used and when should we use a block cipher?

• How should a short password be converted to a 128-bit encryption key?

• What is a side-channel vulnerability?