

MTAT.05.125 Introduction to Theoretical Computer Science

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Practice session 9.

Practise session

1. Define $L = \{0^{2^n} \mid n \in \mathbb{N}\}$, i.e. collection of all strings of zeros whose length is a power of 2. We want to construct TM M that decides L .

We start with implementation-level description. On input w , the machine M :

1. Moves from left to right by crossing out every second zero.
2. If in step 1 the tape contains a single zero – accept.
3. If the tape contains odd number of zeros, which is > 1 , rejects.
4. Return the head to the beginning of the tape and goto 1.

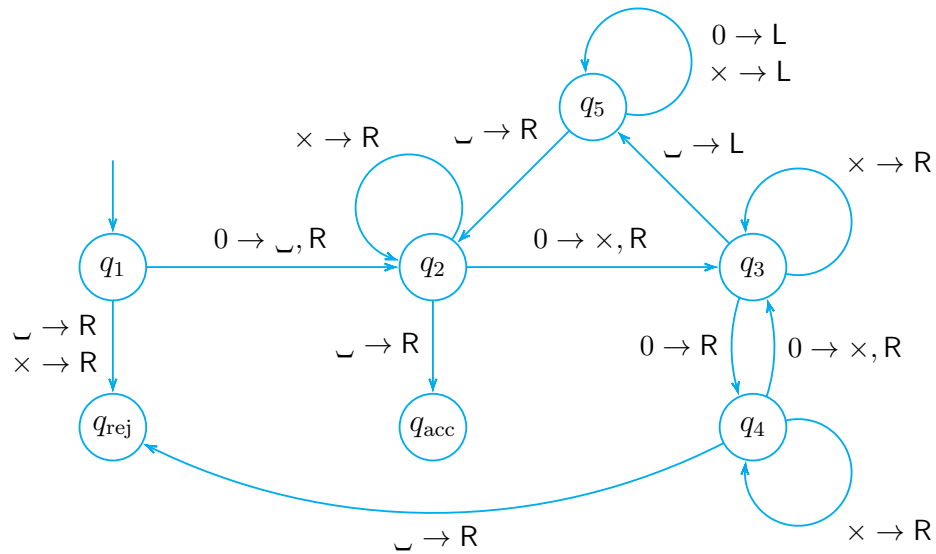
Note that if the number of zeros is 2^n , then on each iteration it decreases by a factor of 2, and after n iterations M accepts.

If it is not a power of 2, then at some iteration condition of the step 3 holds and we reject.

Now we give a formal description. We build TM $M = (Q, \Sigma, \Gamma, \delta, q_1, q_{acc}, q_{rej})$, where

- $Q = \{q_1, q_2, q_3, q_4, q_5, q_{acc}, q_{rej}\}$;
- $\Sigma = \{0\}$;
- $\Gamma = \{0, \times, \sqcup\}$;

- start state is q_1 ;
- accept state is q_{acc} ;
- reject state is q_{rej} ;
- transition function δ is as follows:



Here notation $0 \rightarrow \sqcup, R$ means that by reading 0 M writes \sqcup and moves the head to the right (in transition from q_1 to q_2). Hence, $\delta(q_1, 0) = (q_2, \sqcup, R)$.

Note. M writes \sqcup over the left-most 0 on the tape to mark the left-hand end of the tape.

Let us see the run of M on input 0000:

