

# MTAT.07.017

# Applied Cryptography

Rakenduslik krüptograafia  
Прикладная криптография

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# Topics for This Week

Cryptographic key management

Public key certificate (as defined in X.509)

PKI: Public key infrastructure

# Unsecured Communication



msg = "The Message"

→ → → → → → msg → → → → → →



Problem: anyone can read this message

# Symmetric Encryption



$enc = \text{encrypt}(msg, K)$

$\rightarrow \rightarrow \rightarrow \rightarrow \rightarrow \rightarrow enc \rightarrow \rightarrow \rightarrow \rightarrow \rightarrow \rightarrow$

$msg = \text{decrypt}(enc, K)$



Problem: how to exchange keys?

# Hybrid Encryption



$enc1 = \text{encrypt}(msg, K)$

$enc2 = \text{encrypt}(K, K+)$

$\rightarrow \rightarrow \rightarrow \rightarrow enc1, enc2 \rightarrow \rightarrow \rightarrow \rightarrow$

$K = \text{decrypt}(enc2, K-)$

$msg = \text{decrypt}(enc1, K)$



Problem: how to get  $K+$ ?

# Key Exchange



→ → → (requests  $K^+$ ) → → → →  
← ← ← ← ← ←  $K^+$  ← ← ← ← ← ←

Computes enc1, enc2

→ → → → enc1, enc2 → → → →



Problem: man-in-the-middle attack

# Man-in-the-Middle Attack

Adam believes protocol works like this:

**A** → **B** : Message

**B** → **A** : Response

What is actually happening:

**A** → **M** : Message

**M** → **B** : **AnotherMessage**

**B** → **M** : Reponse

**M** → **A** : **AnotherResponse**



# Man-in-the-Middle Attack

**A** → **M** : request ("PublicKey")

**M** → **B** : request ("PublicKey")

**B** → **M** : **K**+ (**B**)

**M** → **A** : **K**+ (**M**)



**Adam** and **Bob** both believe that **Adam** has received **Bob's** public key





# Man-in-the-Middle Attack

**A** :  $\text{enc1} = \text{encrypt}(\text{msg}, \mathbf{K})$   
**A** :  $\text{enc2} = \text{encrypt}(\mathbf{K}, \mathbf{K} + (\mathbf{M}))$   
**A**  $\rightarrow$  **M** :  $\text{enc1}, \text{enc2}$   
**M** :  $\mathbf{K} = \text{decrypt}(\text{enc2}, \mathbf{K} - (\mathbf{M}))$   
**M** :  $\text{msg} = \text{decrypt}(\text{enc}, \mathbf{K})$   
**M** :  $\text{enc3} = \text{encrypt}(\text{msg}, \mathbf{K})$   
**M** :  $\text{enc4} = \text{encrypt}(\mathbf{K}, \mathbf{K} + (\mathbf{B}))$   
**M**  $\rightarrow$  **B** :  $\text{enc3}, \text{enc4}$



# Key Management

Most problematic part of cryptography?

## Short-term keys

- Session keys – usually generated automatically

## Long-term keys

- Signing and encryption – generated by user

# Key Management Problems

Key physical access control

Key backup/retrieval

Secret sharing schemes

Key destruction

Lifetime of public keys used to validate signatures

# Key Management Models

Offline (PGP, X.509)

Online (S/MIME, SIP, SSH, SSL/TLS)

Ad-hoc

# Offline Key Management

## Hierarchical

Relying on trusted authorities

Example: X.509

## Distributed

Relying on “web of trust”

Example: PGP

# Online Key Management

## On-demand distribution

Public key is sent on request

Example: SSL/TLS, S/MIME

## Key continuity

Regular key exchange initiated

Example: SSH

# Public Key Certificate

Binds public key with key owner's identity

Defines public key validity period

Issued and signed by trusted third party

# Public Key Certificate

```
-- Key-to-identity binding
-- See RFC 2459 for details
TBSCertificate ::= SEQUENCE {
    -- Some fields skipped
validity          Validity,
subject          Name,
subjectPKInfo   SubjectPublicKeyInfo
    -- Some fields skipped
}
```



# Public Key Certificate

```
-- TBSCertificate and authority signature  
-- See RFC 2459 for details  
-- http://tools.ietf.org/html/rfc2459
```

```
Certificate ::= SEQUENCE {  
    tbsCertificate    TBSCertificate,  
    signatureAlg     AlgorithmIdentifier,  
    signatureValue   BIT STRING  
}
```

# X.509 Certificate

Version

Serial Number

Validity Period

Subject DN

Public Key

Extensions

Issuer DN

Issuer Signature

# DN: Distinguished Name

Yet another notation for unique identifiers

Used in LDAP and related protocols

Example:

```
CN=John Doe, OU=Infernal IT, O=Evil Inc., C=US
```

See also: RFC 4514 (<http://tools.ietf.org/html/rfc4514>)

# Requesting a Certificate



```
subj = "CN=John Doe, OU=..."  
{ K+, K- } = generateKeys()  
tbsReq = { subj, K+ }  
reqSig = sign(tbsReq, K-)  
certReq = { tbsReq, reqSig }
```

Certification request `certReq` is sent to CA

# Generating a Certificate

```
result = verifySignature(certReq)
validity = { notBefore, notAfter }
tbsCert = { validity, subj, K+ }
certSig = sign(tbsCert, KCA-)
cert = { tbsCert, certSig }
```

Certificate `cert` is sent back to client

# Problems Obtaining a Certificate

How client locates CA?

How client gets CA's public key?

How client sends own public key to CA?

How CA verifies client's data?

How CA delivers certificate to client?

EUGridPMA: <http://www.eugridpma.org/>

# Questions?

# See You Next Week

No homework for this week (:

Next lab session:

Friday 2011-04-01 **08:30** EET @ Liivi 2 - 205